

## Problem Set Assignment 1 - BNF

The goal of this project is to illustrate how a BNF grammar may be organized and improved to produce a practical or relevant programming language. A "start" symbol, sets of tokens, non-terminal symbols, and productions make up BNF grammar. There will be demonstration and testing of the grammar rules for many tiny languages with different syntax and restrictions.

### Problem 1: Laughter

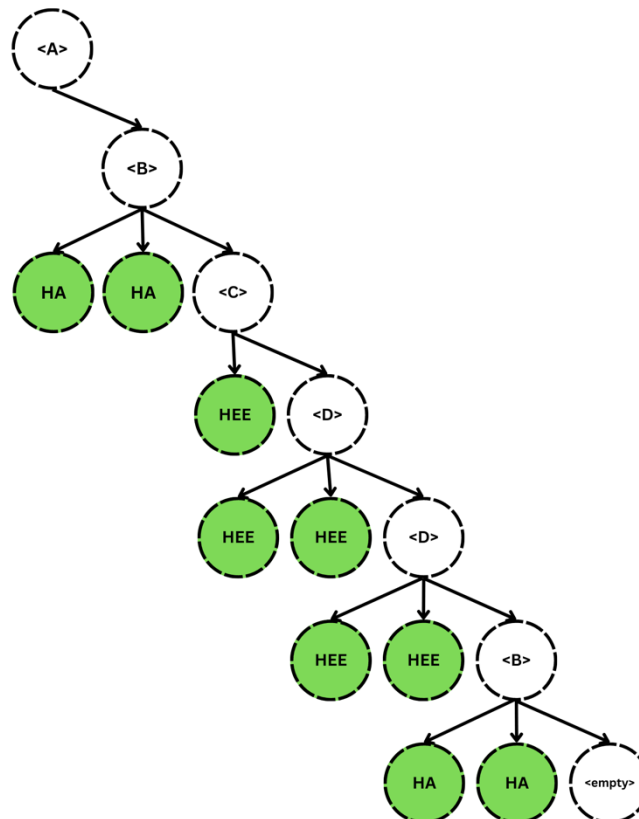
$\langle A \rangle ::= \langle B \rangle \mid \langle C \rangle$

$\langle B \rangle ::= \text{HA HA } \langle C \rangle \mid \text{HA HA } \langle B \rangle \mid \text{HA HA } \langle \text{empty} \rangle$

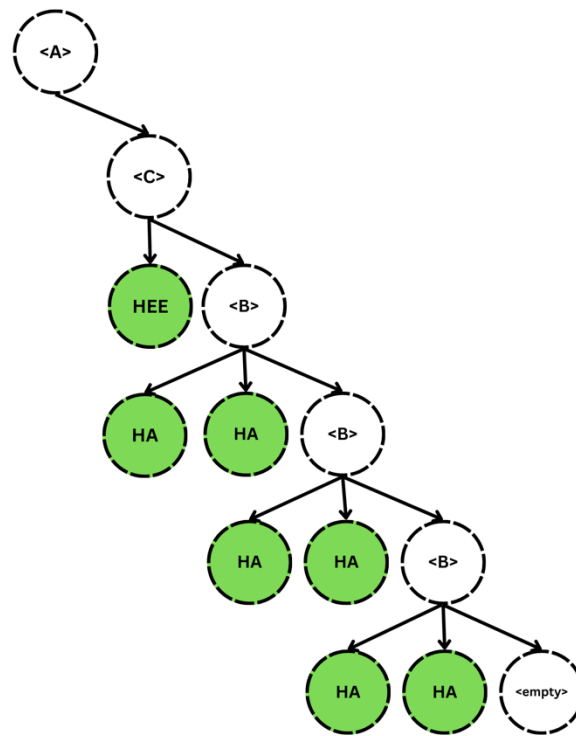
$\langle C \rangle ::= \text{HEE } \langle B \rangle \mid \text{HEE } \langle D \rangle \mid \text{HEE } \langle \text{empty} \rangle$

$\langle D \rangle ::= \text{HEE HEE } \langle B \rangle \mid \text{HEE HEE } \langle D \rangle \mid \text{HEE HEE } \langle \text{empty} \rangle$

- Parse Tree for HA HA HEE HEE HEE HEE HEE HA HA



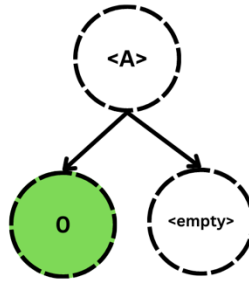
- Parse Tree for HEE HA HA HA HA HA HA



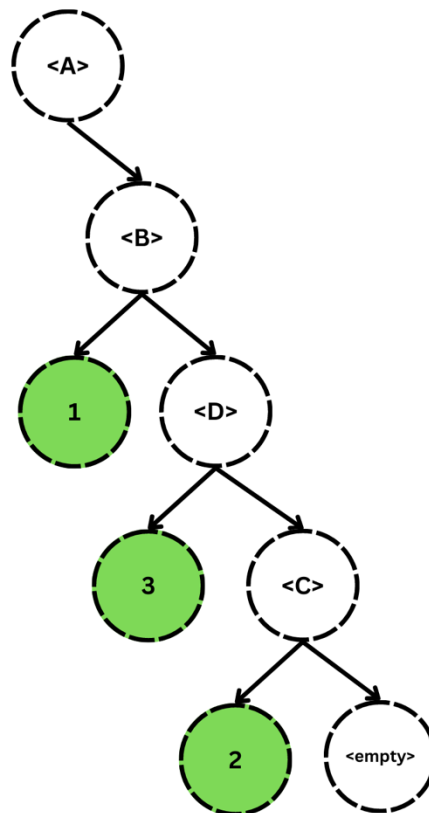
## Problem 2: SQN (Special Quaternary Numbers)

$\langle A \rangle ::= 0 \mid \langle B \rangle \mid \langle C \rangle \mid \langle D \rangle \mid \langle \text{empty} \rangle$   
 $\langle B \rangle ::= 1 \langle A \rangle \mid 1 \langle C \rangle \mid 1 \langle D \rangle \mid \langle \text{empty} \rangle$   
 $\langle C \rangle ::= 2 \langle A \rangle \mid 2 \langle B \rangle \mid 2 \langle D \rangle \mid \langle \text{empty} \rangle$   
 $\langle D \rangle ::= 3 \langle A \rangle \mid 3 \langle B \rangle \mid 3 \langle C \rangle \mid \langle \text{empty} \rangle$

- Parse Tree for 0



- Parse Tree for 132



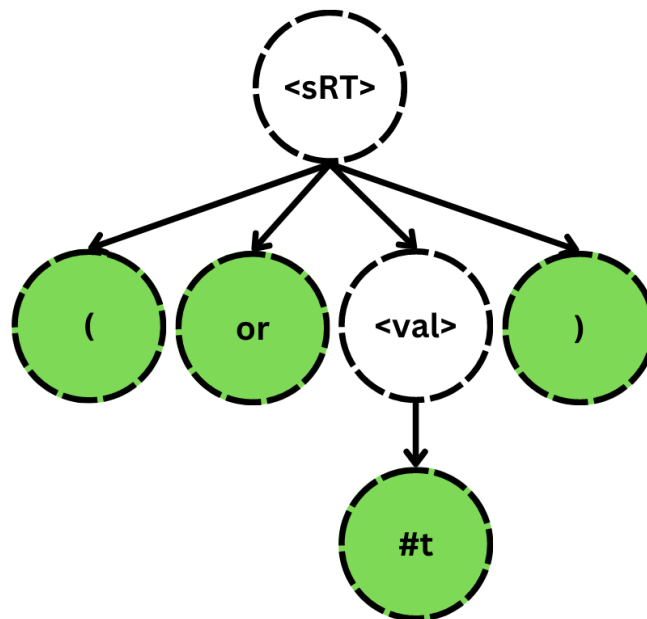
### Problem 3: BXR

```

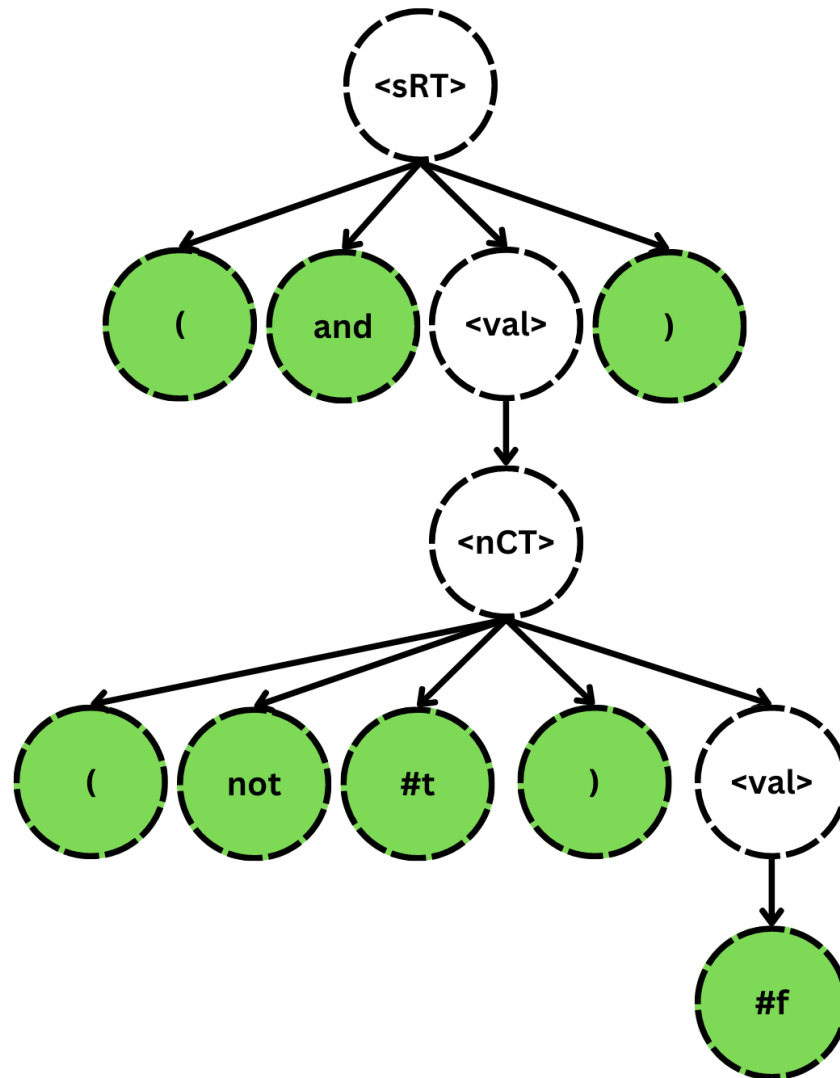
<sRT> ::= ( and <val> ) | ( or <val> ) | <nST> | <oST>
<val> ::= #t <val> | #f <val> | <nCT> | <oCT> | #t | #f
<nCT> ::= ( not #t ) <sRT> | ( not #f ) <sRT> | ( not #t )
<val> | ( not #f ) <val>
<oCT> ::= ( and <val> ) | ( or <val> )
<nST> ::= ( not #t ) | ( not #f )
<oST> ::= ( and #t ) | ( or #t ) | ( and #f ) | ( or #f )

```

- Parse Tree for ( or #t )



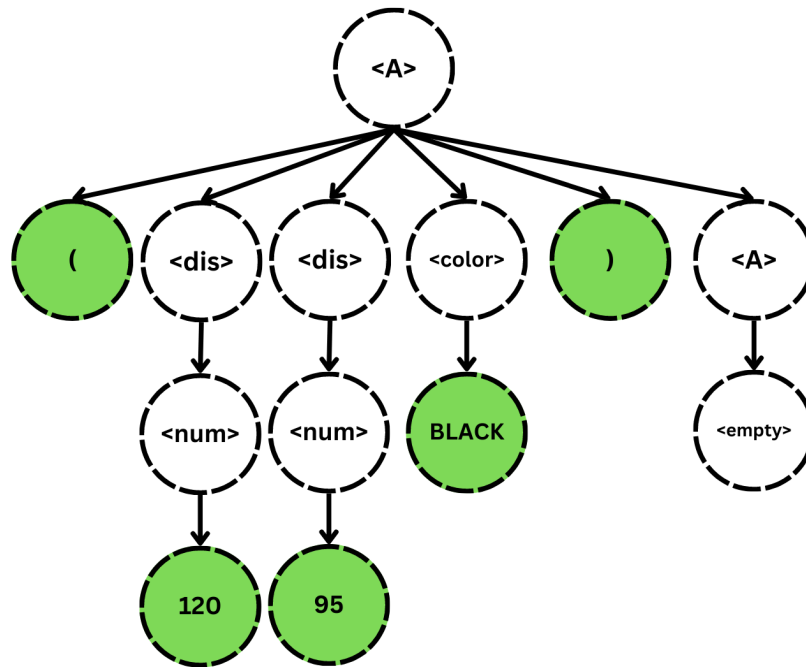
- Parse Tree for ( and ( not #t ) #f )



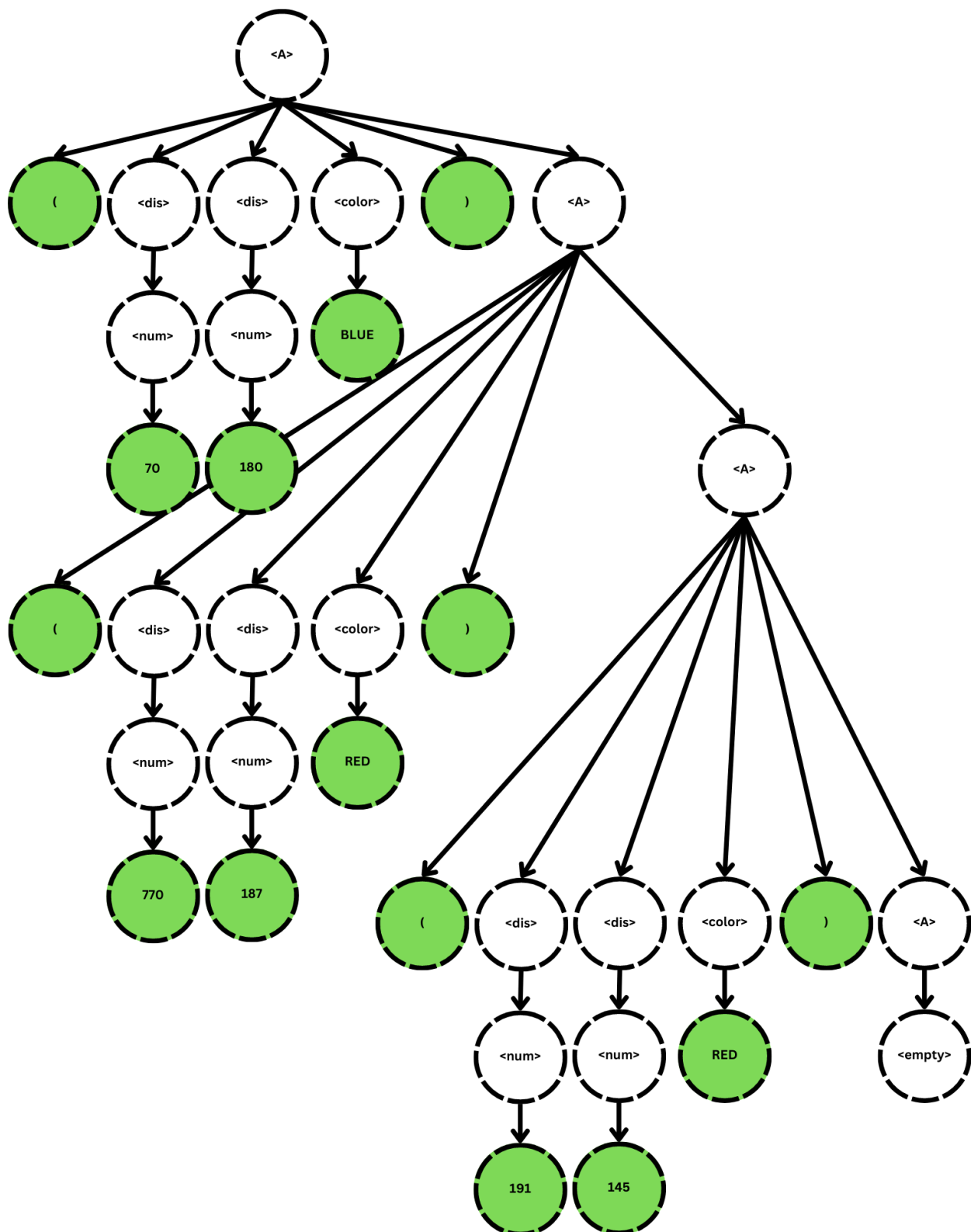
**Problem 4: LSS (Line Segment Sequences)**

$\langle A \rangle ::= ( \langle \text{dis} \rangle \langle \text{ang} \rangle \langle \text{color} \rangle ) \langle A \rangle \mid \langle \text{empty} \rangle$   
 $\langle \text{dis} \rangle ::= \langle \text{num} \rangle$   
 $\langle \text{ang} \rangle ::= \langle \text{num} \rangle$   
 $\langle \text{color} \rangle ::= \text{RED} \mid \text{BLACK} \mid \text{BLUE}$

- Parse Tree for ( 120 95 BLACK )



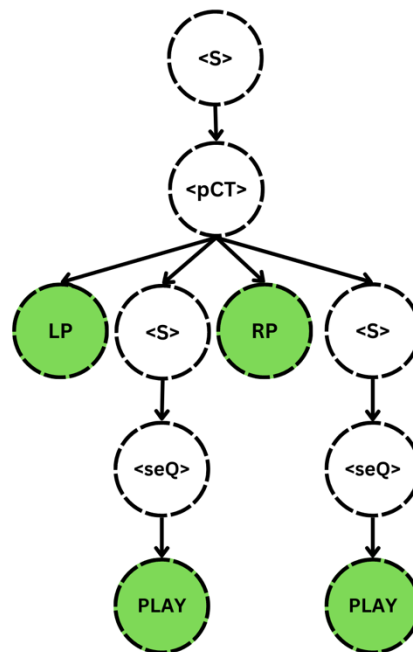
- Parse Tree for ( 70 180 BLUE ) ( 770 187 RED ) ( 191 145 RED )



## Problem 5: M-Lines

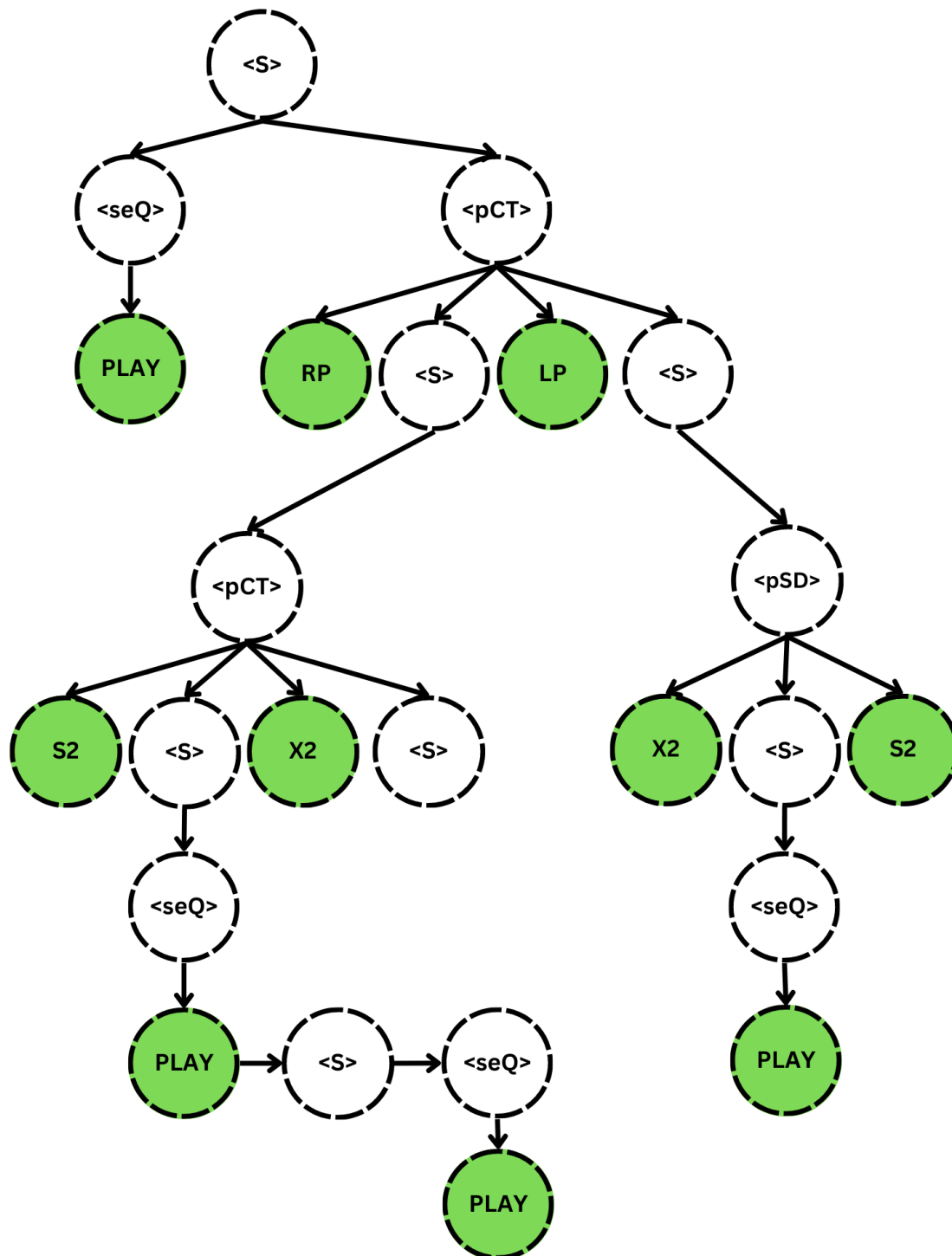
$\langle S \rangle ::= \langle \text{seQ} \rangle \mid \langle \text{pST} \rangle \mid \langle \text{pSD} \rangle \mid \langle \text{pCT} \rangle$   
 $\langle \text{seQ} \rangle ::= \text{PLAY} \mid \text{REST} \mid \text{PLAY} \langle S \rangle \mid \text{REST} \langle S \rangle$   
 $\langle \text{pCT} \rangle ::= \text{RP} \langle S \rangle \text{LP} \langle S \rangle \mid \text{LP} \langle S \rangle \text{RP} \langle S \rangle \mid \text{S2} \langle S \rangle \text{X2} \langle S \rangle \mid \text{X2} \langle S \rangle \text{S2}$   
 $\langle S \rangle \mid \text{S3} \langle S \rangle \text{X3} \langle S \rangle \mid \text{X3} \langle S \rangle \text{S3} \langle S \rangle$   
 $\langle \text{pSD} \rangle ::= \text{RP} \langle S \rangle \text{LP} \mid \text{LP} \langle S \rangle \text{RP} \mid \text{S2} \langle S \rangle \text{X2} \mid \text{X2} \langle S \rangle \text{S2} \mid \text{S3} \langle S \rangle \text{X3} \mid \text{X3} \langle S \rangle \text{S3}$   
 $\langle \text{pST} \rangle ::= \text{RP} \text{LP} \mid \text{LP} \text{RP} \mid \text{S2} \text{X2} \mid \text{X2} \text{S2} \mid \text{S3} \text{X3} \mid \text{X3} \text{S3}$

- Parse Tree for LP PLAY RP PLAY





- Parse Tree for PLAY RP S2 PLAY PLAY X2 LP X2 PLAY S2



## Problem 6: BNF?

Backus-Naur Form, or BNF as it is commonly known, is a system designed to write grammars for programming languages. To understand BNF, one must understand the 4 components that make it up. The first are the “tokens” that are a part of the language according to the definition. Secondly are the symbols that aren’t technically a part of the language but are still essential to its definition. Thirdly, the productions that convert the non-terminal symbols into a string of tokens and nonterminals. Lastly is the start symbol, a nonterminal symbol, which signifies the illustration of the language.