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## *Abstract*

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This assignment is about Backus-Naur Form Grammar (BNF). During this assignment I was asked to compose six different BNF grammars for the different languages given. This also included drawing BNF parse trees and describing them in English in a straightforward, compelling manner. This assignment will expand my knowledge on BNF grammar and parse trees so I can further understand how languages and functions are described.

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## *Problem 1 - Laughter*

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### ***BNF Grammar:***

**Tokens:** { HEE, HA }

**Nonterminals** = { Laughter, Ha, Hee, Hee-Extra }

**Productions** = The following set of rules:

**< Laughter > ::** = < Ha > | < Hee > | < Empty >

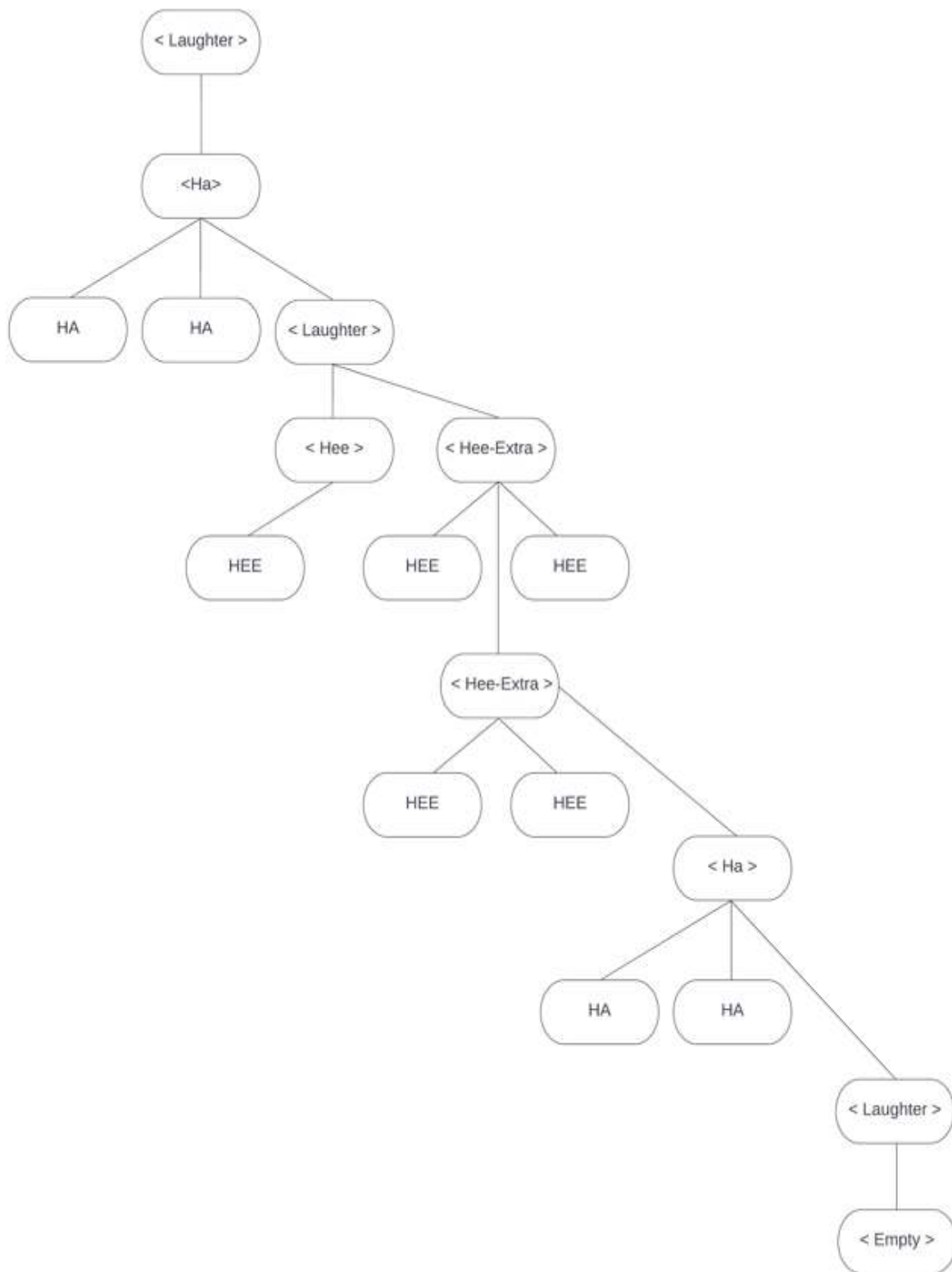
**< Ha > ::** = HA HA <Laughter>

**< Hee > ::** = HEE < Ha > | HEE HEE < Hee-Extra > | < Empty >

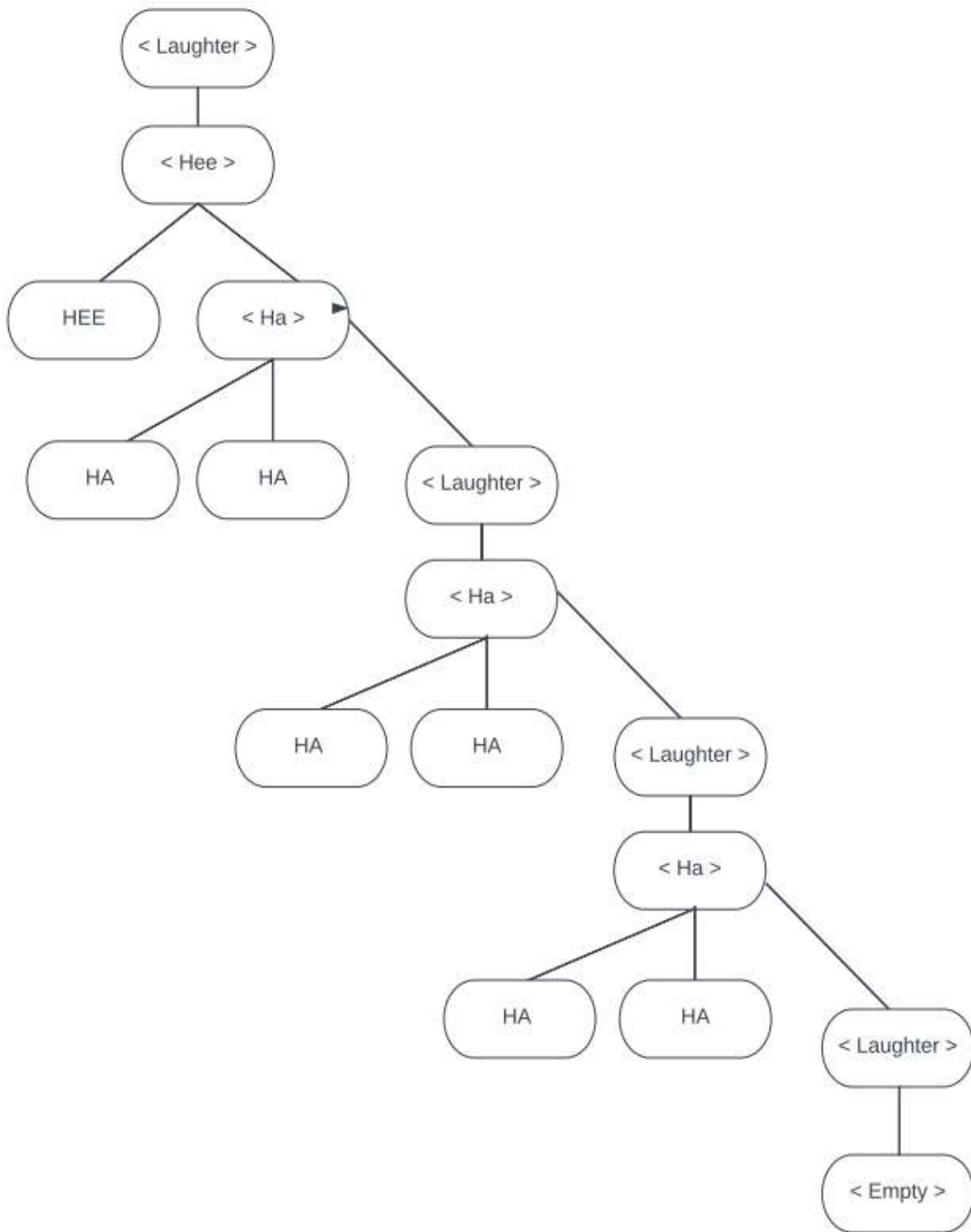
**< Hee-Extra > ::** = Hee Hee < Ha > | < Laughter>

**Start Symbol:** < Laughter >

**Parse Tree Problem - : HA HA HEE HEE HEE HEE HEE HA HA**



## Parse Tree - HEE HA HA HA HA HA HA



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## *Problem 2 - Special Quaternary Numbers*

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**Tokens:** {0, 1, 2, 3}

**Non-Terminals:** { Quart\_Num, Zero, No-Zero, One, No-One, Two, No-Two, Three, No-Three }

**Productions** = The following set of rules

**Quart\_Num ::** = < Zero > | < One > | < Two > | < Three >

**Zero ::** = 0 | < No-Zero >

**No-Zero ::** = < One > | < Two > | < Three > | < Empty >

**One ::** = 1 | < No-One >

**No-One ::** = < Zero > | < Two > | < Three > | < Empty >

**Two ::** = 2 | < No-Two >

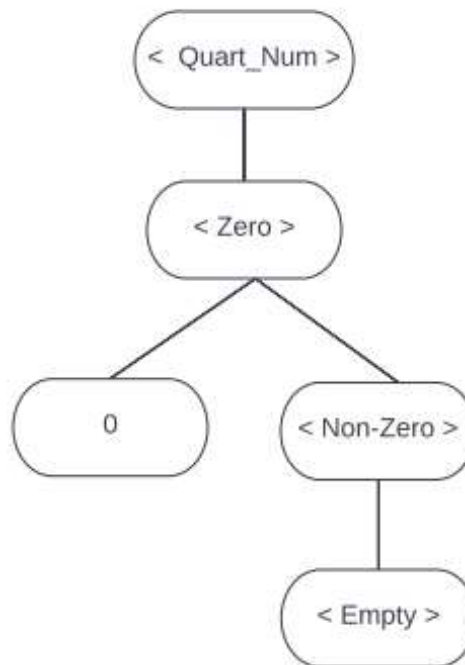
**No-Two ::** = < Zero > | < One > | < Three > | < Empty >

**Three ::** = 3 | < Non-Three >

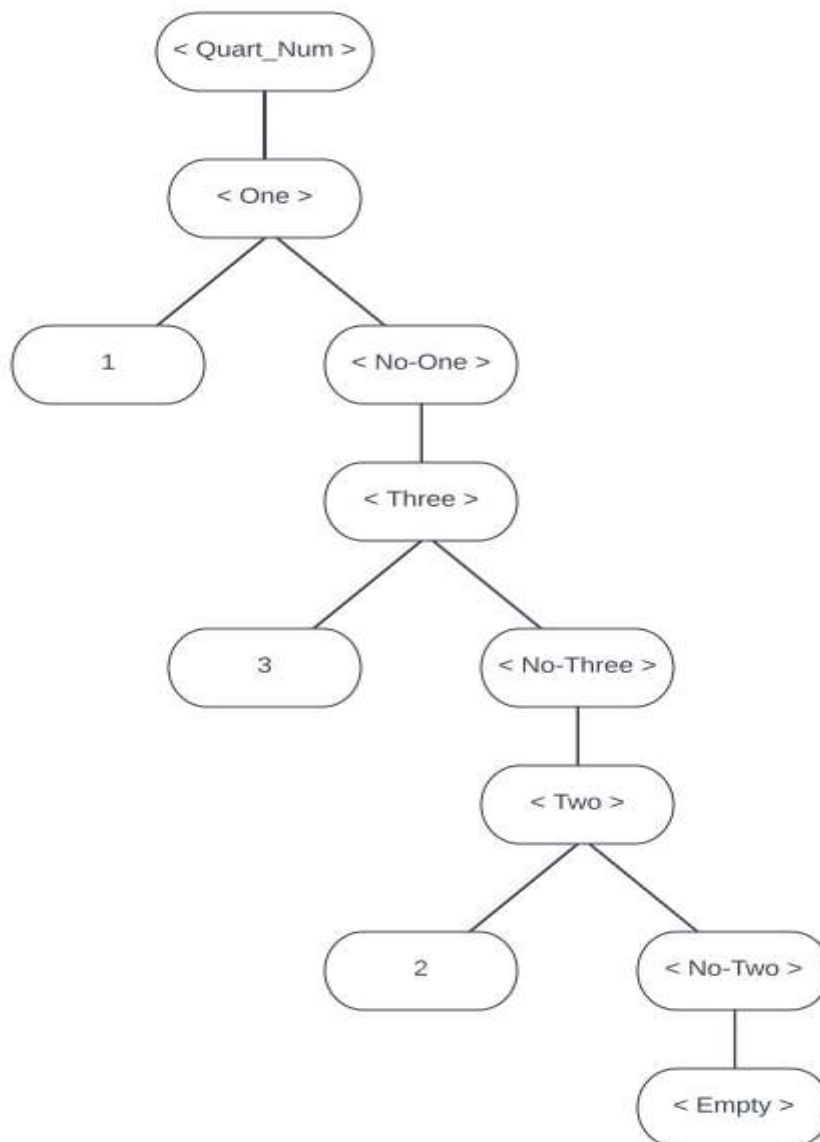
**No-Three ::** = < Zero > | < One > | < Two > | < Empty >

**Starting Symbol:** < Quart\_Num >

**Parse Tree - 0**



## Parse Tree - 132



**Task 4** - Explain why 1223 is not in the language:

The reason 1223 is not in the language is because in the **production** set of rules for state **< Two >** it indicates the only other option is **< No-Two >** which does not allow for another 2 to be adjacent to each other. Therefore 1223 is not in the language.

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## Problem 3 - BXR

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**Non-Terminals:** { BXR, Ops, And, Or, Not, Bool }

**Productions** = The Following Set of Rules

$\langle \mathbf{BXR} \rangle ::= \langle \mathbf{Ops} \rangle \mid \langle \mathbf{Bool} \rangle \mid \text{Empty}$

$\langle \mathbf{Ops} \rangle ::= \langle \mathbf{And} \rangle \mid \langle \mathbf{Or} \rangle \mid \langle \mathbf{Not} \rangle$

$\langle \mathbf{Not} \rangle ::= (\text{not} \langle \mathbf{Bool} \rangle)$

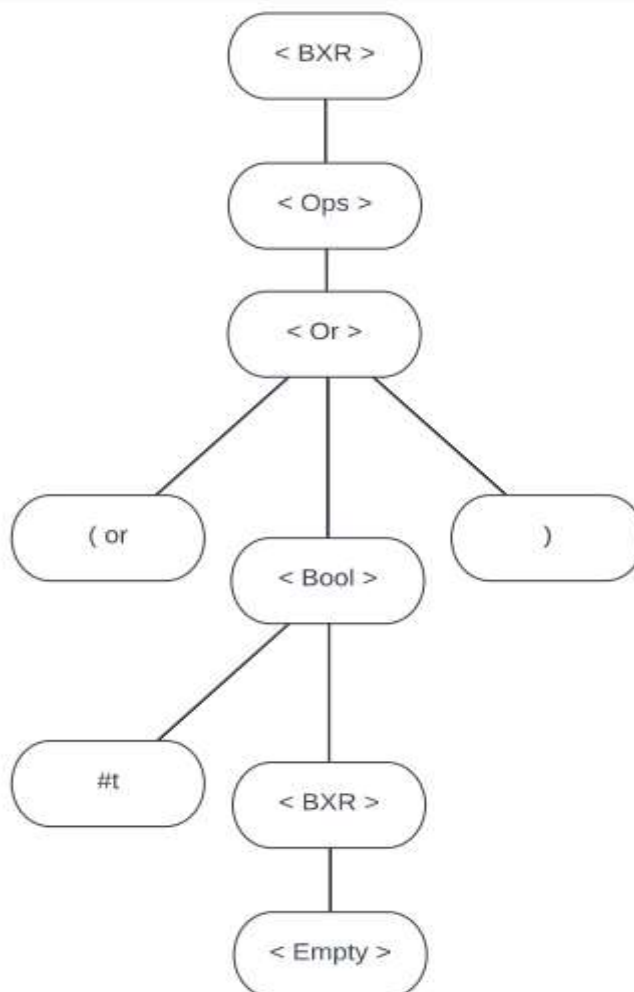
$\langle \mathbf{And} \rangle ::= (\text{and} \langle \mathbf{Bool} \rangle) \mid (\text{and})$

$\langle \mathbf{Or} \rangle ::= (\text{or} \langle \mathbf{Bool} \rangle) \mid (\text{or})$

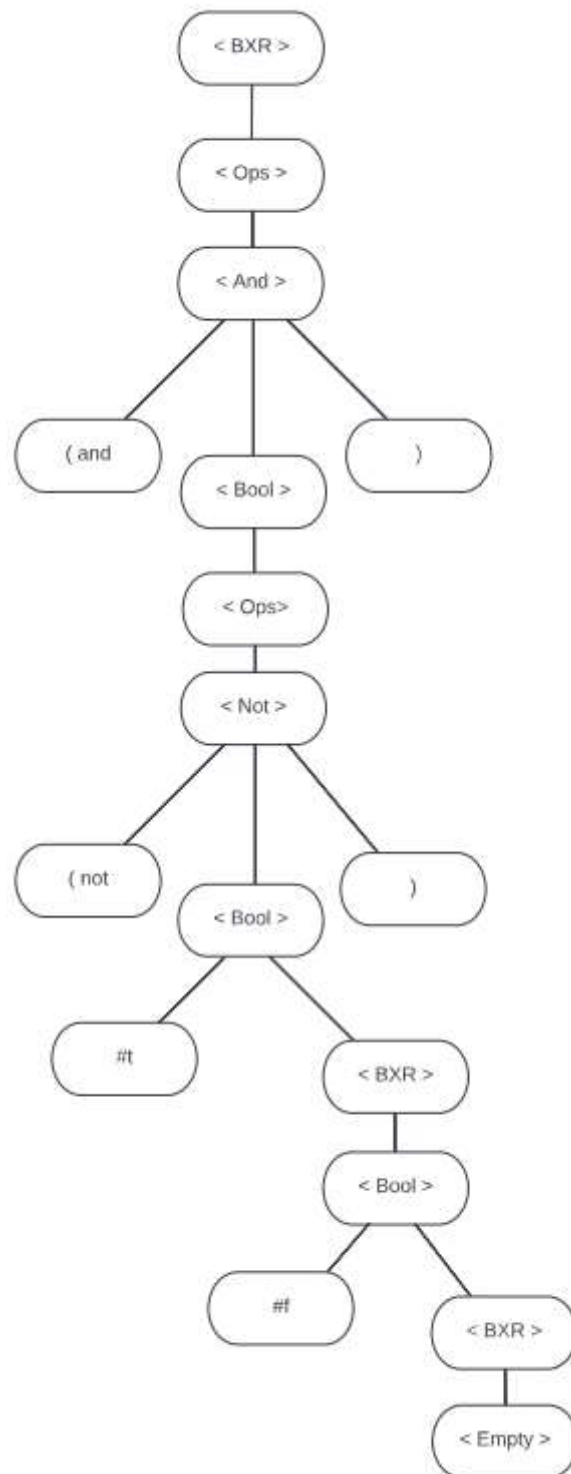
$\langle \mathbf{Bool} \rangle ::= \#t \mid \#f \mid \langle \mathbf{BXR} \rangle \mid \langle \mathbf{Ops} \rangle$

**Starting Symbol:**  $\langle \mathbf{BRX} \rangle$

**Parse Tree - ( or #t )**



## Parse Tree - ( and ( not #t ) #f )



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## Problem 4 - LSS ( Line Segment Sequences )

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**Non-Terminals:** { LSS, Segment, Length, Angle, Color, Empty }

**Productions** = The Following Set of Rules

$\langle \text{LSS} \rangle ::= \langle \text{LSS} \rangle | \langle \text{Segment} \rangle | \langle \text{Empty} \rangle$

$\langle \text{Segment} \rangle ::= ( \langle \text{Length} \rangle \langle \text{Angle} \rangle \langle \text{Color} \rangle )$

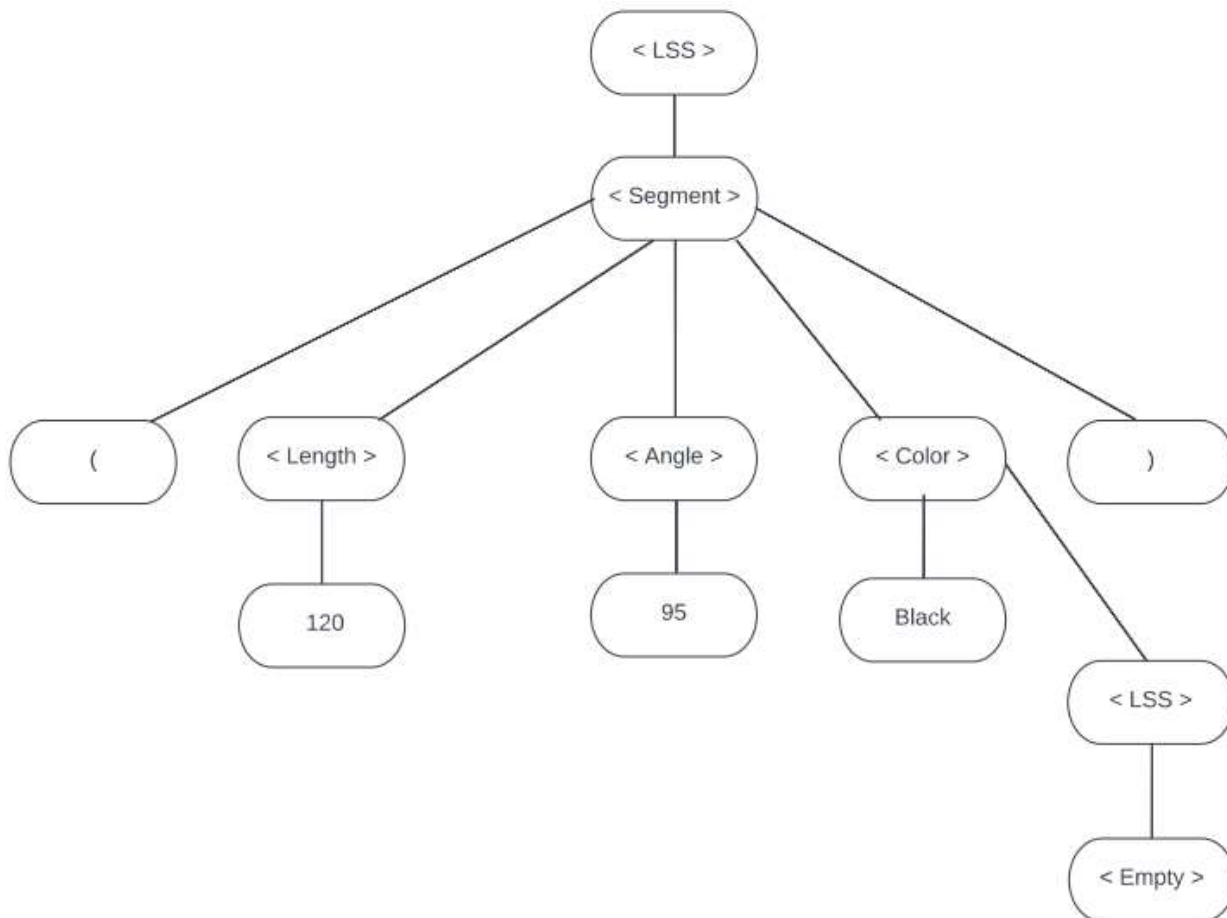
$\langle \text{Length} \rangle ::= \text{int}$

$\langle \text{Angle} \rangle ::= \text{int}$

$\langle \text{Color} \rangle ::= \text{BLACK} | \text{RED} | \text{BLUE} | \langle \text{LSS} \rangle$

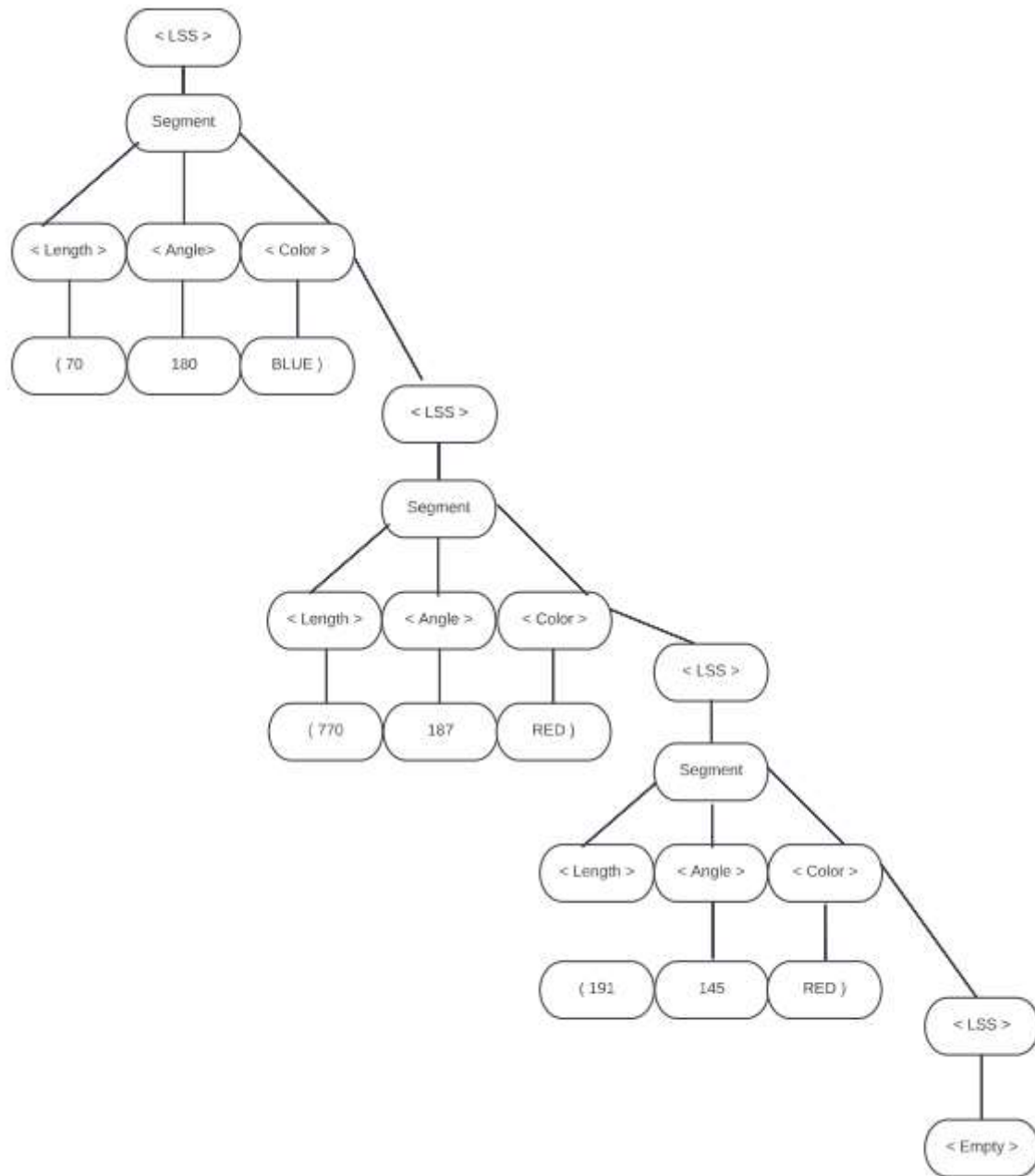
**Starting Symbol:**  $\langle \text{LSS} \rangle$

**Parse Tree - ( 120 95 BLACK )**





# Parse Tree - ( 70 180 BLUE ) ( 770 187 RED ) ( 191 145 RED )



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## Problem 5 - M-Lines

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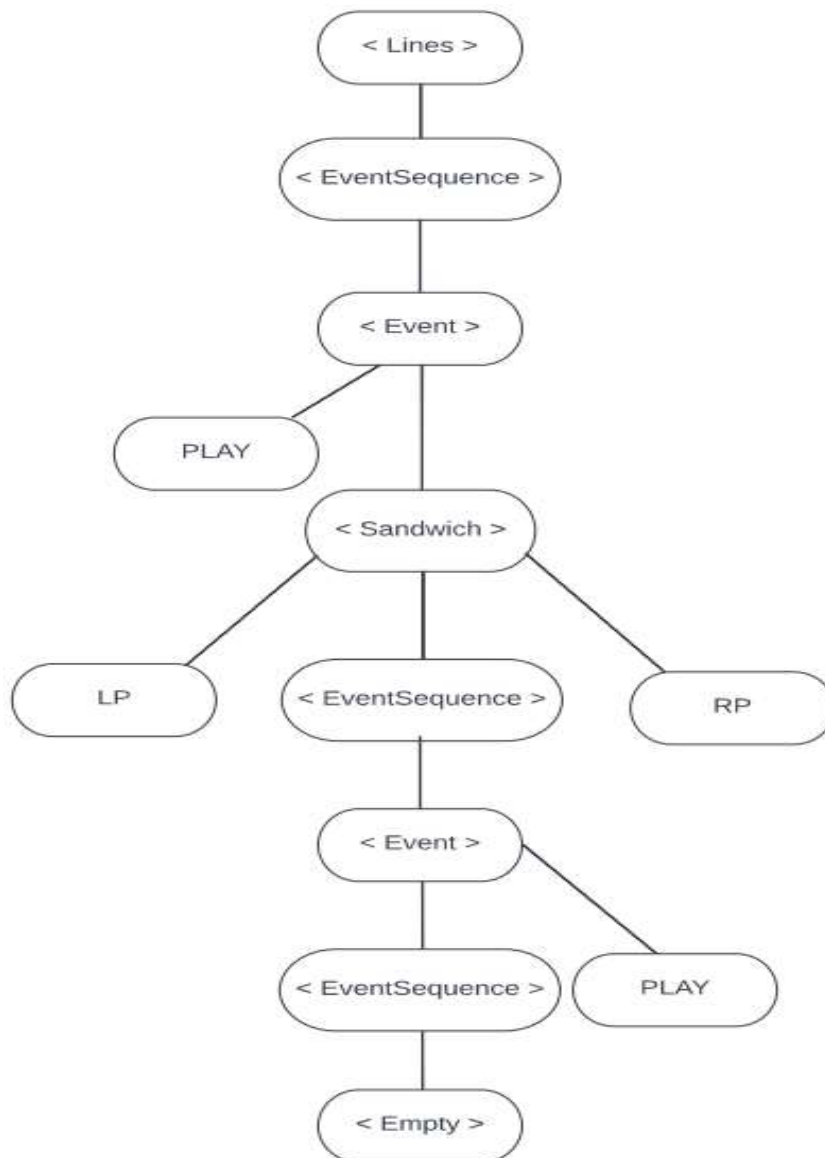
**Non-Terminals:** { Lines, EventSequence, Event, Sandwich }

**Productions** = The Following Set of Rules

$\langle \text{Lines} \rangle ::= \langle \text{EventSequence} \rangle \mid \langle \text{Empty} \rangle$   
 $\langle \text{EventSequence} \rangle ::= \langle \text{Event} \rangle \mid \langle \text{Sequence} \rangle \mid \langle \text{Event} \rangle \langle \text{Sequence} \rangle \mid$   
 $\quad \langle \text{Sequence} \rangle \langle \text{Event} \rangle \mid \langle \text{Empty} \rangle$   
 $\langle \text{Event} \rangle ::= \text{PLAY} \langle \text{EventSequence} \rangle \mid \text{REST} \langle \text{EventSequence} \rangle$   
 $\langle \text{Sandwich} \rangle ::= \text{RP} \langle \text{EventSequence} \rangle \text{LP} \mid \text{LP} \langle \text{EventSequence} \rangle \text{RP} \mid$   
 $\quad \text{S2} \langle \text{EventSequence} \rangle \text{X2} \mid \text{X2} \langle \text{EventSequence} \rangle \text{S2} \mid$   
 $\quad \text{S3} \langle \text{EventSequence} \rangle \text{X3} \mid \text{X3} \langle \text{EventSequence} \rangle \text{S3}$

**Starting Symbol:**  $\langle \text{Lines} \rangle$

**Parse Tree - LP PLAY RP PLAY**



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## *Problem 6 - BNF?*

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BNF Otherwise known as Backus-Naur form is a type of notation that is issued to describe the syntax of languages such as programming languages and instruction sets. It consists of a set of terminal symbols, set of non terminal symbols and a set of production rules of the form. These symbols are then used to describe the instructions or execution of a program often done using parse trees.